

March 23–25, 2026 Inria Workshop 2026

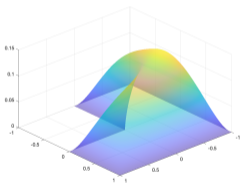
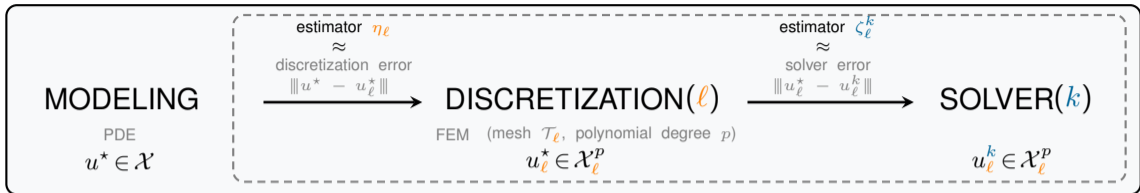
AFEM with optimally preconditioned GMRES

Ani Miraçi, [Laboratoire Jacques-Louis Lions](#)

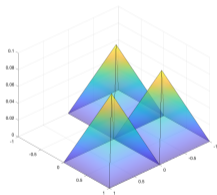
joint work with Thomas Führer, Paula Hilbert, Dirk Praetorius



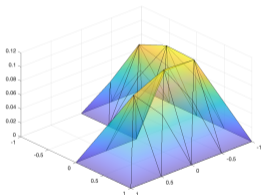
[link to slides](#)



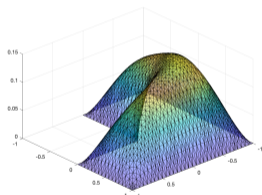
solution u^*



$\ell = 0$
 $k = 1$



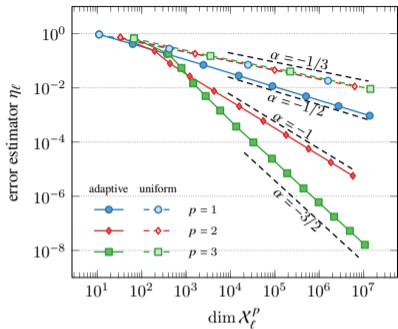
$\ell = 1$
 $k = 3$



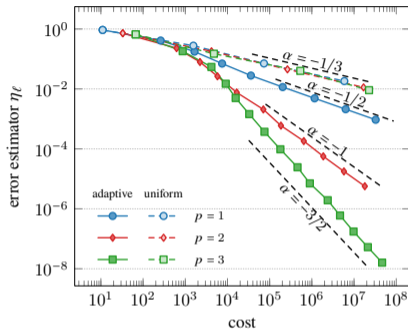
$\ell = 8$
 $k = 0$

1. **Reliability** $\|u^* - u_\ell^k\| \lesssim \eta_\ell + \zeta_\ell^k$
2. **Unconditional convergence** $u_\ell^k \rightarrow u^*$
3. **Optimal rates of convergence**





Optimal convergence $\alpha = p/d$ wrt **dofs** and **cost**



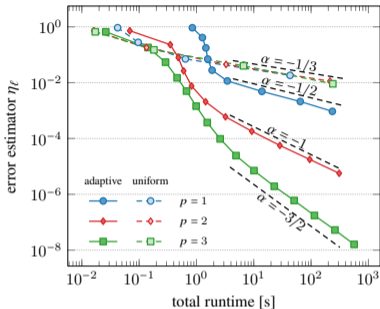
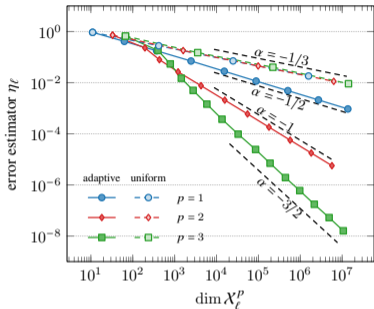
$$\sup_{\ell \in \mathbb{N}_0} (\#\mathcal{T}_\ell)^\alpha \eta_\ell < \infty$$



$$\sup_{\ell \in \mathbb{N}_0} \left(\sum_{\ell'=0}^{\ell} \sum_{k'=0}^{k[\ell']} \#\mathcal{T}_{\ell'} \right)^\alpha \eta_\ell < \infty$$

-  Binev, Dahmen, DeVore: *Numer. Math.*, 97 (2004)
-  Stevenson: *Found. Comput. Math.*, 7 (2007)
-  Cascón, Kreuzer, Nochetto, Siebert: *SIAM J. Numer. Anal.*, 46 (2008)
-  Gantner, Haberl, Praetorius, Schimanko: *Math. Comp.*, 90 (2021)

Towards optimal complexity



Optimal complexity of AFEM requires each of its modules to be realized in linear complexity:

- SOLVE is **critical**
- ESTIMATE ✓
- MARK (Stevenson 2007, Pfeiler-Praetorius 2020 for minimal cardinality marking) ✓
- REFINE (Binev-Dahmen-DeVore 2004, Stevenson 2008) ✓

Optimal algebraic solver

Model problem

Consider the **non-symmetric linear elliptic PDE**

$$-\operatorname{div}(\mathbf{K}\nabla u^*) + \mathbf{b} \cdot \nabla u^* + cu^* = f - \operatorname{div} \mathbf{f} \quad \text{in } \Omega \quad \text{with } u^* = 0 \quad \text{on } \partial\Omega$$

Weak formulation: search for $u^* \in \mathcal{X} := H_0^1(\Omega)$ solution of

$$b(u^*, v) = \langle f, v \rangle_{L^2(\Omega)} + \langle \mathbf{f}, \nabla v \rangle_{L^2(\Omega)} =: F(v) \quad \text{for all } v \in \mathcal{X}$$

Notation: $a(v, w) := \langle \mathbf{K} \nabla v, \nabla w \rangle_{L^2(\Omega)} \quad \|v\| := a(v, v)^{1/2}$

$$b(v, w) := a(v, w) + \langle \mathbf{b} \cdot \nabla v, w \rangle_{L^2(\Omega)} + \langle cv, w \rangle_{L^2(\Omega)} \quad \text{for all } v, w \in \mathcal{X}$$

Lax-Milgram framework \implies well-posed problem

$$|b(v, w)| \leq C_{\text{bnd}} \|v\| \|w\| \quad \text{and} \quad b(v, v) \geq C_{\text{ell}} \|v\|^2 \quad \text{for all } v, w \in \mathcal{X}$$

FEM discretization

Functional writing

- \mathcal{T}_ℓ simplicial mesh of Ω and $p \geq 1$ polynomial degree
- finite element space $\mathcal{X}_\ell^p := \{v_\ell \in H_0^1(\Omega) : v_\ell|_T \in \mathbb{P}_p(T) \ \forall T \in \mathcal{T}_\ell\} \subset \mathcal{X}$
- search for $u_\ell^* \in \mathcal{X}_\ell^p$ solution of

$$b(u_\ell^*, v_\ell) = F(v_\ell) \quad \forall v_\ell \in \mathcal{X}_\ell^p$$

Discrete writing:

- $N_\ell := \dim \mathcal{X}_\ell^p$
- $\varphi_{\ell,1}, \dots, \varphi_{\ell,N_\ell}$ basis of \mathcal{X}_ℓ^p
- $(B_\ell)_{j,k} := b(\varphi_{\ell,k}, \varphi_{\ell,j})$ $(d_\ell)_k := F(\varphi_{\ell,k})$ $(A_\ell)_{j,k} := a(\varphi_{\ell,k}, \varphi_{\ell,j})$
- search for $\mathbf{x}_\ell^* \in \mathbb{R}^{N_\ell}$ solution of

$$B_\ell \mathbf{x}_\ell^* = \mathbf{d}_\ell$$

Link:

$$u_\ell^* := \sum_{j=1}^{N_\ell} (\mathbf{x}_\ell^*)_j \varphi_{\ell,j}$$

Optimal additive Schwarz

Preconditioner requirements:

- symmetric and positive definite preconditioner $P_\ell \approx B_\ell^{-1}$

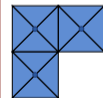
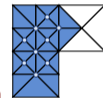
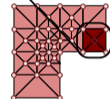
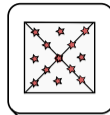
$$\|\mathbf{x}_\ell\|_{P_\ell^{-1}} := (\mathbf{x}_\ell, \mathbf{x}_\ell)_{P_\ell^{-1}}^{1/2} := (P_\ell^{-1} \mathbf{x}_\ell, \mathbf{x}_\ell)_2^{1/2}$$

- linear complexity of applying P_ℓ
- discrete-functional equivalence

$$C_b^{-1} \|v_\ell\| \leq \|\mathbf{x}_\ell\|_{P_\ell^{-1}} \leq C_b \|v_\ell\| \quad \text{for all } v_\ell = \sum_{j=1}^{N_\ell} (\mathbf{x}_\ell)_j \varphi_{\ell,j} \in \mathcal{X}_\ell^p$$

Subspace correction via additive Schwarz:

$$P_\ell^{AS} := \underbrace{I_0^1 \underbrace{(A_0^1)^{-1}}_{\substack{\text{lowest-order} \\ \text{global solve}}} (I_0^1)^T + \sum_{\ell'=1}^{\ell-1} I_{\ell'}^1 \underbrace{(D_{\ell'}^+)^{-1}}_{\substack{\text{lowest-order} \\ \text{local solves}}} (I_{\ell'}^1)^T}_{h\text{-robustness}} + \underbrace{\sum_{z \in \mathcal{V}_\ell} I_{\ell,z}^p \underbrace{(A_{\ell,z}^p)^{-1}}_{\substack{\text{high-order} \\ \text{patch-wise solves}}} (I_{\ell,z}^p)^T}_{p\text{-robustness}}$$



 Chen, Nochetto, Xu: *Numer. Math.*, 120 (2012)

 Schöberl, Melenk, Pechstein, Zaglmayr: *IMA J. Numer. Anal.*, 28 (2008)

 Hilbert, Miraçi, Praetorius: arXiv: 2601.14911 (2026)

PGMRES with restart

input: system matrix B_ℓ , preconditioner P_ℓ , right-hand side d_ℓ , initial guess x_ℓ^0 , $\Lambda : \mathbb{R}^{N_\ell} \rightarrow \mathbb{R}_{\geq 0}$, $k_{\max} \geq 1$

initialize $r_\ell^0 := d_\ell - B_\ell x_\ell^0$ $s_\ell^0 := P_\ell r_\ell^0$ $\tilde{v}^1 := r_\ell^0 / \|s_\ell^0\|_{P_\ell^{-1}} = r_\ell^0 / \langle s_\ell^0, r_\ell^0 \rangle_2^{1/2}$ $v^1 := P_\ell \tilde{v}^1$ $k := -1$

for each $\tilde{k} = 0, 1, 2, \dots, k_{\max}$ **repeat** (residual minimization loop)

$\tilde{w} := B_\ell v^{\tilde{k}}$ $w := P_\ell \tilde{w}$ $k \leftarrow k + 1$ $R \leftarrow \lfloor k/k_{\max} \rfloor$

for $j = 1, 2, \dots, \tilde{k}$, **repeat** (Arnoldi process loop)

$H_{j,\tilde{k}} := \langle w, v^j \rangle_{P_\ell^{-1}} = \langle \tilde{w}, v^j \rangle_2$ $\tilde{w} := \tilde{w} - H_{j,\tilde{k}} \tilde{v}^j$ $w := w - H_{j,\tilde{k}} v^j$

end

$H_{\tilde{k}+1,\tilde{k}} := \|w\|_{P_\ell^{-1}} = \langle \tilde{w}, w \rangle_2^{1/2}$ $y := \arg \min_{z \in \mathbb{R}^{\tilde{k}}} \| \|s_\ell^{Rk_{\max}}\|_{P_\ell^{-1}} e^1 - Hz \|_2$

$x_\ell^k := x_\ell^{Rk_{\max}} + \sum_{j=1}^{\tilde{k}} y_j v^j$ $r_\ell^k := d_\ell - B_\ell x_\ell^k$ $s_\ell^k := P_\ell r_\ell^k$ $\|s_\ell^k\|_{P_\ell^{-1}} = \langle r_\ell^k, s_\ell^k \rangle_2^{1/2}$

if $\tilde{k} \neq k_{\max}$ **then** $\tilde{v}^{\tilde{k}+1} := \tilde{w} / H_{\tilde{k}+1,\tilde{k}}$, $v^{\tilde{k}+1} := w / H_{\tilde{k}+1,\tilde{k}}$ **else** new basis from r_ℓ^k

until $\|s_\ell^k\|_{P_\ell^{-1}} \leq \Lambda(x_\ell^k)$

set $k[\ell] := k$

output: approximation $x_\ell^{k[\ell]}$, restart index R , preconditioned residual norm $\|s_\ell^{k[\ell]}\|_{P_\ell^{-1}}$

A posteriori estimator and solver contraction

Proposition 1 (h - and p -robust **reliability** and **efficiency** of the a posteriori estimator)

$$C_{\text{ell}} C_b^{-2} \|u_\ell^* - u_\ell^k\| \leq \|s_\ell^k\|_{P_\ell^{-1}} =: \zeta_\ell^k \leq C_{\text{bnd}} C_b^2 \|u_\ell^* - u_\ell^k\|$$

Proposition 2 (h - and p -robust **contraction**)

$$\|s_\ell^{k+1}\|_{P_\ell^{-1}} \leq q_{\text{alg}} \|s_\ell^k\|_{P_\ell^{-1}} \quad 0 < q_{\text{alg}}(C_{\text{bnd}}, C_{\text{ell}}, C_b) < 1$$

Note: result holds for any $k_{\text{max}} \geq 1 \implies$ linear complexity

Challenge: contraction yes, but in preconditioner norm: \implies even if $\|v_{\ell+1}^0\| = \|v_\ell^k\|$
 \implies typically $\|s_{\ell+1}^0\|_{P_{\ell+1}^{-1}} \neq \|s_\ell^k\|_{P_\ell^{-1}}$

 Eisenstat, Elman, Schultz: *SIAM J. Numer. Anal.*, 20 (1983)

 Sarkis, Szyld: *Comput. Methods Appl. Mech. Engrg.*, 196 (2007)

 Schöberl, Melenk, Pechstein, Zanglmayr: *IMA J. Numer. Anal.*, 28 (2008)

 Miraçi, Papež, Vohralík: *SIAM J. Sci. Comput.*, 43 (2021)

 Innerberger, Miraçi, Praetorius, Streitberger: *ESAIM Math. Model. Numer. Anal.*, 58 (2024)

 Führer, Hilbert, Miraçi, Praetorius: (in preparation)

Adaptive mesh refinement and algebraic solver

AFEM with iterative solver

input: initial mesh \mathcal{T}_0 , initial guess u_0^0 , adaptivity parameters $0 < \theta \leq 1$, λ_{alg} , C_{alg} , $S := 0$

for each $\ell = 0, 1, 2, \dots$ **repeat**

(mesh-refinement loop)

SOLVE & ESTIMATE

(algebraic solver loop)

compute u_ℓ^k and $\|s_\ell^k\|_{P_\ell^{-1}}$ from u_ℓ^0 via PGMRES with $\Lambda(x_\ell^k) := \lambda_{\text{alg}} \eta_\ell(u_\ell^k)$ for $k \leq k[\ell]$

ADAPTIVE PARAMETER CONTROL

- **if** $S > C_{\text{alg}} (\eta_\ell(u_\ell^k) + \|s_\ell^k\|_{P_\ell^{-1}})^{-1}$ **then** $C_{\text{alg}} \leftarrow 2 C_{\text{alg}}$ $\lambda_{\text{alg}} \leftarrow \lambda_{\text{alg}}/2$
- $S \leftarrow S + (\eta_\ell(u_\ell^k) + \|s_\ell^k\|_{P_\ell^{-1}})^{-1}$

MARK select quasi-minimal $\mathcal{M}_\ell \subseteq \mathcal{T}_\ell$ such that $\theta \sum_{T \in \mathcal{T}_\ell} \eta_\ell(T, u_\ell^k)^2 \leq \sum_{T \in \mathcal{M}_\ell} \eta_\ell(T, u_\ell^k)^2$

REFINE $\mathcal{T}_{\ell+1} := \text{refine}(\mathcal{T}_\ell, \mathcal{M}_\ell)$

$$u_{\ell+1}^0 := u_\ell^k$$

output: approximations u_ℓ^k and corresponding estimators $\zeta_\ell^k = \|s_\ell^k\|_{P_\ell^{-1}}, \eta_\ell(u_\ell^k)$

Reliability and linear convergence

A posteriori control for each approximation u_ℓ^k

$$\underbrace{\|u^\star - u_\ell^k\|}_{\text{total error}} \leq \underbrace{\|u^\star - u_\ell^\star\|}_{\text{discretization error}} + \underbrace{\|u_\ell^\star - u_\ell^k\|}_{\text{algebraic error}} \stackrel{\text{reliability (discretization)}}{\lesssim} \underbrace{\eta_\ell(u_\ell^\star)}_{\text{discretization estimator}} + \|u_\ell^\star - u_\ell^k\|$$

$$\stackrel{\text{reliability (algebra)}}{\lesssim} \eta_\ell(u_\ell^k) + \underbrace{\|s_\ell^k\|_{P_\ell^{-1}}}_{\text{algebra estimator}} =: \underbrace{H_\ell^k}_{\text{quasi-error}}$$

Define the index set $\mathcal{Q} := \{(\ell, k) \in \mathbb{N}^2 : u_\ell^k \text{ is computed by algorithm}\}$
 with ordering $|\ell, k| := \#\{(\ell', k') \in \mathcal{Q} : u_{\ell'}^{k'} \text{ computed earlier than } u_\ell^k\}$

Definition (full R-linear convergence of the quasi-error)

$$H_\ell^k \leq C_{\text{lin}} q_{\text{lin}}^{|\ell, k| - |\ell', k'|} H_{\ell'}^{k'}, \quad C_{\text{lin}} > 0, \quad 0 < q_{\text{lin}} < 1$$

Contraction independently of the algorithmic step: mesh-refinement or algebraic iteration

Corollary (convergence of the total error)

$$\|u^\star - u_\ell^k\| \lesssim H_\ell^k \lesssim q_{\text{lin}}^{|\ell, k|} H_0^0 \longrightarrow 0 \quad \text{for} \quad |\ell, k| \longrightarrow \infty$$

Conditional full R-linear convergence

Lemma (summability criteria and R-linear convergence) Let $a_\ell \in \mathbb{R}_{\geq 0}$, $\ell \in \mathbb{N}_0$, $\beta, \gamma > 0$. Then (i)–(iii) are pairwise equivalent:

(i) *R-linear convergence*: $a_\ell \leq C_{\text{lin}} q_{\text{lin}}^{\ell-\ell'} a_{\ell'}$ for all $0 \leq \ell' \leq \ell$ $0 < q_{\text{lin}} < 1, C_{\text{lin}} > 1$

(ii) *tail summability*: $\sum_{\ell'=\ell+1}^{\infty} a_{\ell'}^\beta \leq C_\beta a_\ell^\beta$ for all $\ell \in \mathbb{N}_0$ $C_\beta > 0$

(iii) *inverse summability*: $\sum_{\ell'=0}^{\ell-1} a_{\ell'}^{-\gamma} \leq C_\gamma a_\ell^{-\gamma}$ for all $\ell \in \mathbb{N}_0$ $C_\gamma > 0$

If (i)–(iii) hold for $\ell \geq \ell_0 \in \mathbb{N}$, then they hold for all $\ell \in \mathbb{N}_0$

Proposition (conditional tail summability in the levels) Let $\lambda^* := \min\{1, C_{\text{stab}}^{-1} C_{\text{ell}} C_b^{-2}\}$. If there exists $\ell_0 \in \mathbb{N}$ such that $0 < \lambda_{\text{alg}} < \lambda^* \theta^{1/2}$ for all $\ell \geq \ell_0$, then

$$\sum_{\ell'=\ell+1}^{\infty} H_{\ell'}^k \leq C_{\text{tail}} H_\ell^k \quad C_{\text{tail}} \geq 1$$

Unconditional full R-linear convergence

ADAPTIVE PARAMETER CONTROL: $H_\ell^k = \eta_\ell(u_\ell^k) + \|s_\ell^k\|_{P_\ell}^{-1}$ *computable*

- if $S = \sum_{\ell'=0}^{\ell-1} (H_{\ell'}^k)^{-1} > C_{\text{alg}} (H_\ell^k)^{-1}$ then $C_{\text{alg}} \leftarrow 2 C_{\text{alg}}$ $\lambda_{\text{alg}} \leftarrow \lambda_{\text{alg}}/2$
- $S \leftarrow S + (H_\ell^k)^{-1}$

\implies new a-posteriori-steered criterion algorithmically enforces the inverse summability criterion

Lemma (finitely many updates) The input parameters C_{alg} and λ_{alg} are updated only finitely many times.

Theorem (unconditional full R-linear convergence of the quasi-error) Consider arbitrary $0 < \theta \leq 1$, $\lambda > 0$. Then

$$H_\ell^k \leq C_{\text{lin}} q_{\text{lin}}^{|\ell, k| - |\ell', k'|} H_{\ell'}^{k'}, \quad C_{\text{lin}} > 0, \quad 0 < q_{\text{lin}} < 1$$

Full R-linear convergence implies **rates = complexity**

- $\mathfrak{R}(\alpha) := \sup_{(\ell, k) \in \mathcal{Q}} (\#\mathcal{T}_\ell)^\alpha H_\ell^k < \infty$ rate α wrt dofs is **possible**
- $\widehat{\mathfrak{R}}(\alpha) := \sup_{(\ell, k) \in \mathcal{Q}} \left(\sum_{\substack{(\ell', k') \in \mathcal{Q} \\ |\ell', k'| \leq |\ell, k|}} \#\mathcal{T}_{\ell'} \right)^\alpha H_\ell^k < \infty$ rate α wrt costs/overall computational cost is **possible**

Proposition

$$\mathfrak{R}(\alpha) \leq \widehat{\mathfrak{R}}(\alpha) \leq \frac{C_{\text{lin}}}{(1 - q_{\text{lin}}^{1/\alpha})^\alpha} \mathfrak{R}(\alpha)$$

► **Proof:** $\#\mathcal{T}_{\ell'} \leq \mathfrak{R}(\alpha)^{\frac{1}{\alpha}} (H_{\ell'}^{k'})^{-\frac{1}{\alpha}} \quad \forall (\ell', k') \in \mathcal{Q}$, summing and using the **geometric series** :

$$\Rightarrow \sum_{\substack{(\ell', k') \in \mathcal{Q} \\ |\ell', k'| \leq |\ell, k|}} \#\mathcal{T}_{\ell'} \leq \mathfrak{R}(\alpha)^{\frac{1}{\alpha}} \sum_{\substack{(\ell', k') \in \mathcal{Q} \\ |\ell', k'| \leq |\ell, k|}} (H_{\ell'}^{k'})^{-\frac{1}{\alpha}} \leq \mathfrak{R}(\alpha)^{\frac{1}{\alpha}} C_{\text{lin}}^{\frac{1}{\alpha}} \left(\sum_{\substack{(\ell', k') \in \mathcal{Q} \\ |\ell', k'| \leq |\ell, k|}} q_{\text{lin}}^{(|\ell, k| - |\ell', k'|)/\alpha} \right) (H_\ell^k)^{-\frac{1}{\alpha}}$$

Optimal complexity

We introduce the notion of *approximation class* :

$$\|u^*\|_{\mathbb{A}_\alpha} := \sup_{N \geq \#\mathcal{T}_0} N^\alpha \left[\min_{\#\mathcal{T}_{\text{opt}} \leq N} \eta_{\text{opt}} \right]$$

$$\|u^*\|_{\mathbb{A}_\alpha} < \infty \iff u^* \text{ can be approximated with rate } \alpha \text{ wrt dofs}$$

Theorem (optimal convergence wrt to overall computational cost)

Let $\alpha > 0$ such that $\|u^*\|_{\mathbb{A}_\alpha} < \infty$. Suppose $0 < \theta < 1$ and $\lambda_{\text{alg}} > 0$ **sufficiently small**

$$\implies \|u^*\|_{\mathbb{A}_\alpha} \lesssim \sup_{(\ell, k) \in \mathcal{Q}} \left(\sum_{\substack{(\ell', k') \in \mathcal{Q} \\ |\ell', k'| \leq |\ell, k|}} \#\mathcal{T}_{\ell'} \right)^\alpha H_\ell^k \lesssim \max \{ \|u^*\|_{\mathbb{A}_\alpha}, H_0^0 \}$$

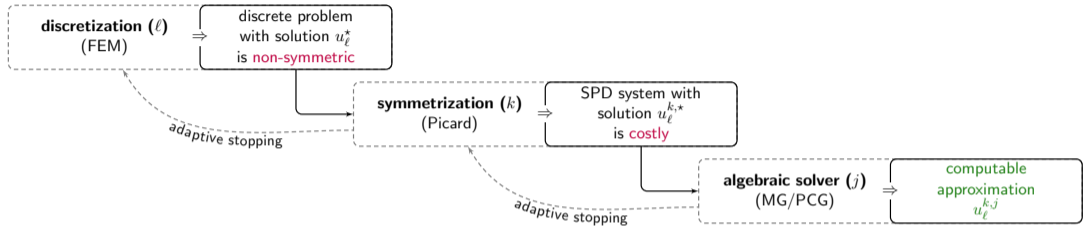
if u^* can be approximated with rate α wrt dofs

then AFEM with optimal algebraic solver approximates u^* with the rate α wrt cost.

 Gantner, Haberl, Praetorius, Schimanko: *Math. Comp.*, 90 (2021)

 Führer, Hilbert, Miraçi, Praetorius: (in preparation)

Adaptive approach with nested solvers



⇒ three **nested** loops: mesh-refinement (ℓ) \mapsto symmetrization (k) \mapsto algebra (j)

⇒ each of the solvers (algebra and symmetrization) is **contractive**

 Congreve, Wihler: *J. Comput. Appl. Math.*, 311 (2017)

 Heid, Praetorius, Wihler: *Comput. Methods Appl. Math.*, 21 (2021)

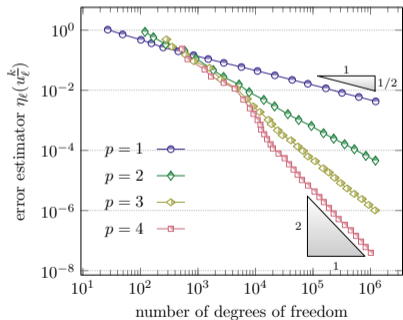
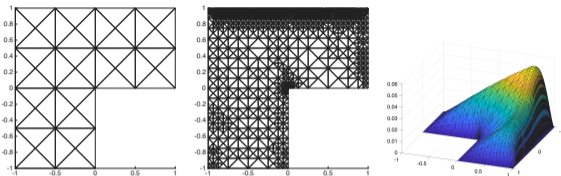
 Haberl, Praetorius, Schimanko, Vohralík: *Numer. Math.*, 147 (2021)

 Brunner, Innerberger, Miraçi, Praetorius, Streitberger, Heid: *IMA J. Numer. Anal.*, 44 (2024)

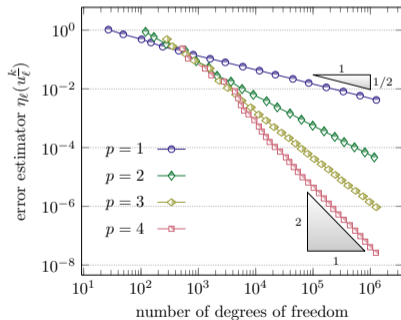
Numerical experiments

Rate-optimality of AFEM

$$-\Delta u^* + \begin{pmatrix} 1 \\ 25 \end{pmatrix} \cdot \nabla u^* = 1 \quad \text{in } \Omega.$$

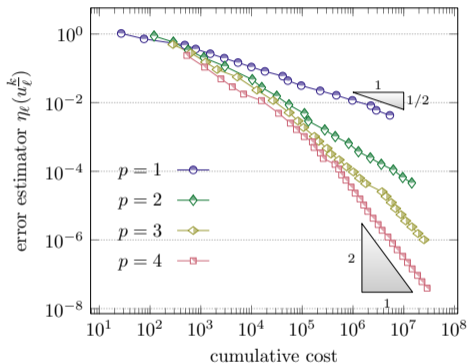


$k_{\max} = 5$

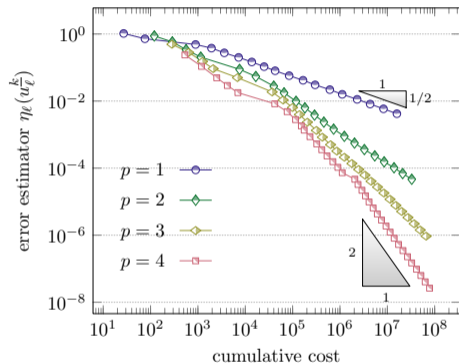


$k_{\max} = 1$

Cost-optimality of AFEM

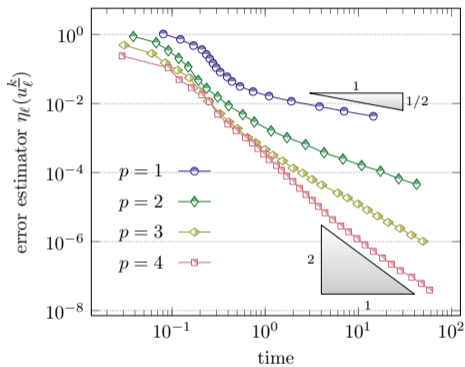


$k_{\max} = 5$

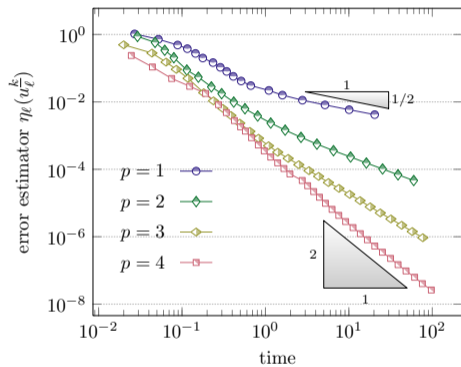


$k_{\max} = 1$

Optimal complexity of AFEM

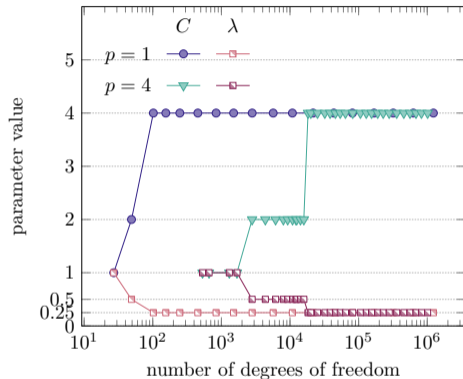


$$k_{\max} = 5$$

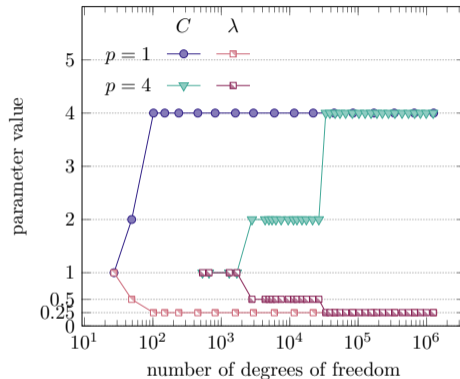


$$k_{\max} = 1$$

Parameter control

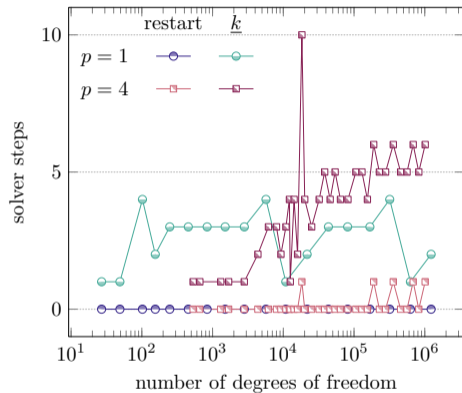


$k_{\max} = 5$

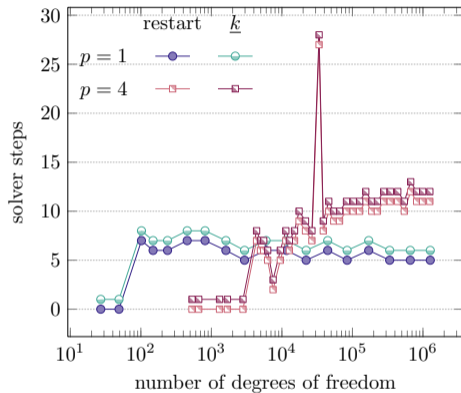


$k_{\max} = 1$

PGMRES steps

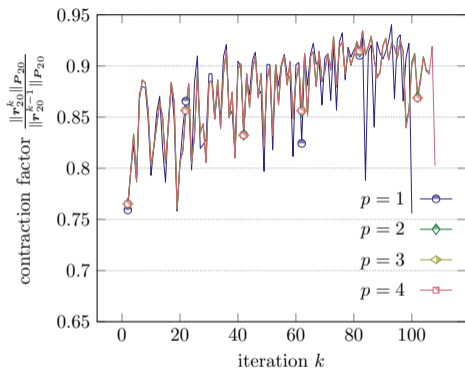


$$k_{\max} = 5$$

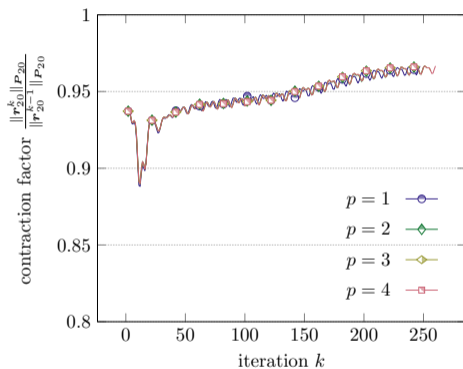


$$k_{\max} = 1$$

PGMRES robustness

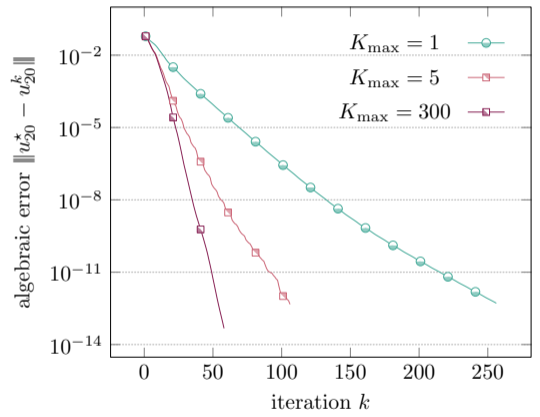
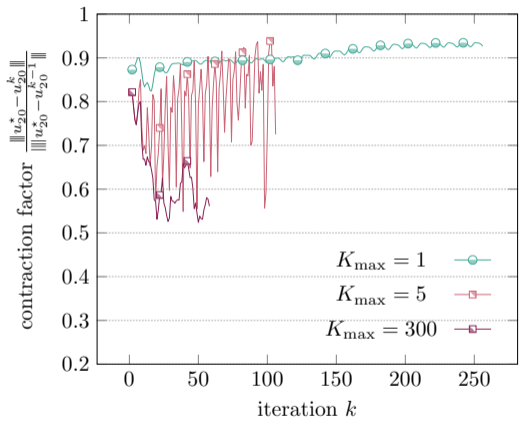


$k_{\max} = 5$



$k_{\max} = 1$

PGMRES restart and energy error study



Take home messages

- 1 analysis of AFEM for non-symmetric problems *without nested solvers*
- 2 *linear complexity* and *optimality* of the preconditioner is crucial
- 3 PGMRES *contracts* in preconditioned norm even with *restart* after each step
- 4 use *nested* iterations and *termination* of solver balancing the different error components
- 5 *reliability* via a posteriori error estimators is ensured at each algorithmic step
- 6 *new adaptive control* on parameters enforcing algorithmically full R-linear convergence
 - ▶ gives *contraction* regardless of algorithmic step
 - ▶ holds for *arbitrary* adaptivity parameters
 - ▶ provides the equivalence *rates = complexity*
- 7 *optimal complexity* is ensured for sufficiently small parameters

Selected contributions

 Miraçi, Papež, Vohralík

A-posteriori-steered p -robust multigrid with optimal step-sizes and adaptive number of smoothing steps
SIAM J. Sci. Comput., 43, DOI: 10.1137/20M1349503 (2021)

 Innerberger, Miraçi, Praetorius, Streitberger

hp -robust multigrid solver on locally refined meshes for FEM discretizations of symmetric elliptic PDEs
ESAIM Math. Model. Numer. Anal., 58, DOI: 10.1051/m2an/2023104 (2024)

 Bringmann, Miraçi, Praetorius

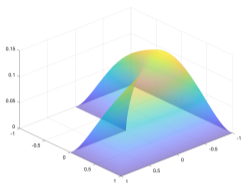
Chapter Four - Iterative solvers in adaptive FEM: Adaptivity yields quasi-optimal computational runtime

Advances in Applied Mechanics, Elsevier, 59, DOI:
<https://doi.org/10.1016/bs.aams.2024.08.002> (2024)

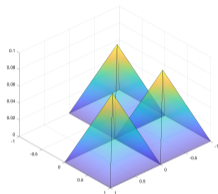
 Führer, Hilbert, Miraçi, Praetorius

AFEM with optimally preconditioned GMRES yields optimal complexity
(in preparation)

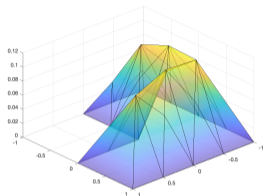
Thank you for your attention!



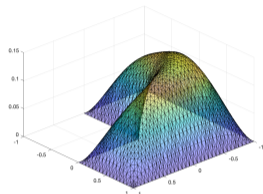
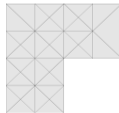
solution u^*



$\ell = 0$
 $k = 1$



$\ell = 1$
 $k = 3$



$\ell = 8$
 $k = 0$

